

A few definitions first:

**Definition 1** A graph  $G = (V, E)$  is called “dense” if for each pair of non-adjacent vertices  $u$  and  $v$ ,  $d(u) + d(v) \geq n$  where  $n = |V|$  and  $d(\bullet)$  denotes the degree of the vertex  $\bullet$ .

**Definition 2** A “Hamiltonian cycle” on  $G$  is a sequence of vertices  $(v_{i_1} v_{i_2} \dots v_{i_n} v_{i_1})$  such that  $v_{i_l} \neq v_{i_h}$  for all  $l \neq h$  and  $\{v_{i_l}, v_{i_{l+1}}\}$  is an edge of  $G$ .

The problem is: write a program that, given a dense indirect graph  $G = (V; E)$  as input, determines whether  $G$  admits a Hamiltonian cycle on  $G$  and outputs that cycle, if there is one, or outputs ‘N’ if there is none.

## Input

The input file contains several descriptions of graphs (each one ending with a ‘%’), in the form:

```
n1
ui1 uj1
ui2 uj2
...
%
n2
ui1 uj1
ui2 uj2
...
%
```

where  $n_i$  is the number of vertices ( $0 < n_i \leq 256$ ) and  $u_{i_h} u_{i_l}$  are integers between 1 and  $n_i$  indicating that there exists an edge between vertex  $u_{i_h}$  and  $u_{i_l}$

## Output

For each test case, output a line that must contain the sequence of vertices that form a Hamiltonian cycle in the form:

```
ui1 ui2 ui3 ...
```

or containing:

```
N
```

## Sample Input

```
4
1 2
2 3
2 4
3 4
3 1
%
6
1 2
1 3
1 6
3 2
3 4
5 2
5 4
6 5
6 4
%
```

## Sample Output

```
1 2 4 3 1
1 3 2 5 4 6 1
```